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## Contributions

- A novel Schur Complement preconditioner
- A framework to apply different solvers to inner subdomains
- High performance method to solve the pressure projection problem for incompressible flows


## Overview

- Previous efforts in production
- Schur complement method and existing preconditioners
- Our new preconditioner
- Choosing subdomain solvers
- Poisson solver tests and fluid simulation tests
- Limitations and future work


## Production Inspiration

- ILM
- Scanline
- Weta
- Simplifie sims

- But we usually use ad-hoc solutions
- Requires clever artists
- Tweaked for each shot


## Incompressible Euler Equations

$\frac{\partial \mathbf{u}}{\partial t}=-(\mathbf{u} \cdot \nabla) \mathbf{u}-\frac{1}{\rho} \nabla \mathbf{p}+\mathbf{f}$
$\nabla \cdot \mathbf{u}=\mathbf{0}$

## Slow!

## Inspiration

- Lot of prior work in CFD and graphics
- Adaptive fluid simul

A scalable Schur-complement fluids solver for heterogeneous compute platforms

- Fast Poisson solve
- Multi-grid
- domain decomposit

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## Fast Poisson Solver

- Regular domain: n log n
- Boundaries using IOP [Molemaker et al. 2008]
- Large scale gas sims [Henderson 2012]
- Not usable for liquid simulations
- But PCG has poor parallel scalability
- Cores are idling


## Design Requirements

- 1 billion voxels
- Parallelized for multi-core systems
- Use the Fast Poisson solver algorithm
- PIC/FLIP is popular
- Large sims are slow, and need too much memory
- Our target resolutions needed NB-FLIP [Ferstl et al. 2016]


## Schur Complement Decomposition



## - Subdomain

$\square$ Boundary set

$$
\left(\begin{array}{ccccc}
A_{11} & & & & A_{1 B} \\
& A_{22} & & & A_{2 B} \\
& & \ddots & & \vdots \\
A_{1 B}^{T} & A_{2 B}^{T} & \cdots & A_{n B}^{T} & A_{n B} \\
A_{B B}
\end{array}\right)\left(\begin{array}{c}
x_{1} \\
x_{2} \\
\vdots \\
x_{n} \\
x_{B}
\end{array}\right)=\left(\begin{array}{c}
b_{1} \\
b_{2} \\
\vdots \\
b_{n} \\
b_{B}
\end{array}\right)
$$

## Schur Complement Matrix

$$
\left(\begin{array}{ccccc}
A_{11} & & & & A_{1 B} \\
& A_{22} & & & A_{2 B} \\
& & \ddots & & \vdots \\
& & & A_{n n} & A_{n B} \\
A_{1 B}^{T} & A_{2 B}^{T} & \cdots & A_{n B}^{T} & A_{B B}
\end{array}\right)\left(\begin{array}{c}
x_{1} \\
x_{2} \\
\vdots \\
x_{n} \\
x_{B}
\end{array}\right)=\left(\begin{array}{c}
b_{1} \\
b_{2} \\
\vdots \\
b_{n} \\
b_{B}
\end{array}\right)
$$

$$
\begin{aligned}
& \left\{\begin{array}{c}
A_{11} x_{1}+A_{1 B} x_{B}=b_{1}, \\
A_{22} x_{2}+A_{2 B} x_{B}=b_{2}, \\
\ldots \ldots ., \\
A_{n n} x_{n}+A_{n B} x_{B}=b_{n} .
\end{array}\right. \\
& x_{i}=A_{i i}^{-1}\left(b_{i}-A_{i B} x_{B}\right)
\end{aligned}
$$

## Schur Complement Matrix

$$
\begin{aligned}
& A_{1 B}^{T} x_{1}+A_{2 B}^{T} x_{2}+\ldots+A_{n B}^{T} x_{n}+A_{B B} x_{B}=b_{B} \\
& S x_{B}=b \\
& S=A_{B B}-\sum_{i=1}^{n} A_{i B}^{T} A_{i i}^{-1} A_{i B} \\
& b=b_{B}-\sum_{i=1}^{n} A_{i B}^{T} A_{i i}^{-1} b_{i}
\end{aligned}
$$

## Schur Complement Solver

$$
\begin{aligned}
& S x_{B}=b \\
& S=A_{B B}-\sum_{i=1}^{n} A_{i B}^{T} A_{i i}^{-1} A_{i B}
\end{aligned}
$$

- $S$ is also symmetric positive definite
- We don't need to form S explicitly
- We use PCG to solve this
- We only need to multiply a vector with S


## Schur Complement Solver


$A_{i i} w_{i}=A_{i B} w_{B}$

- Actually solve the subdomain equation

$$
S w_{B}=A_{B B} w_{B}-\sum_{i=1}^{n} A_{i B}^{T} w_{i}
$$

## Subdomain Independence

$$
\begin{aligned}
& S x_{B}=b \\
& S=A_{B B}-\sum_{i=1}^{n} A_{i B}^{T} A_{i i}^{-1} A_{i B} \\
& b=b_{B}-\sum_{i=1}^{n} A_{i B}^{T} A_{i i}^{-1} b_{i}
\end{aligned}
$$

- Aii subdomains are all independent
- Subdomains can be solved in parallel


## Block Jacobi Preconditioner



## Cross Points Subdomain

- The cross points are disconnected with subdomains
- Treat them as another subdomain
- Modified system

$$
\left(\begin{array}{cccccc}
A_{11} & & & & A_{1 B} \\
& \ddots & & & & \vdots \\
& & A_{n n} & & A_{n B} \\
& & & A_{W W} & A_{W B} \\
A_{1 B}^{T} & \cdots & A_{n B}^{T} & A_{W B}^{T} & A_{B B}
\end{array}\right)\left(\begin{array}{c}
x_{1} \\
\vdots \\
x_{n} \\
x_{W} \\
x_{B}
\end{array}\right)=\left(\begin{array}{c}
b_{1} \\
\vdots \\
b_{n} \\
b_{n} \\
b_{W} \\
b_{B}
\end{array}\right)
$$

## The 3D View



- Wirebasket and cross points are disconnected from the subdomains
- Cross points are also disconnected from face sets


## Our New Preconditioner



View the subdomain as solid

$$
\left(\begin{array}{cc}
A_{W W} & A_{W B^{\prime}} \\
A_{W B^{\prime}}^{T} & N_{B^{\prime} B^{\prime}}
\end{array}\right)\binom{v_{W}}{v_{B^{\prime}}}=\binom{0}{r_{B^{\prime}}}
$$



## The "Aha" Moment

- This is still too slow to solve
- Edge sets are disconnected

- Formulate the preconditioner as
a Schur complement problem
- And we can parallelize the preconditioner too!


## Our New Preconditioner

$$
\left(\begin{array}{ll}
A_{W W} & A_{W B^{\prime}} \\
A_{W B^{\prime}}^{T} & N_{B^{\prime} B^{\prime}}
\end{array}\right)\binom{v_{W}}{v_{B^{\prime}}}=\binom{0}{r_{B^{\prime}}}
$$

Make another Schur system

$$
\left(\begin{array}{ccccc}
N_{F_{1} F_{1}} & & & & A_{F_{1} W^{\prime}} \\
& \ddots & & & \\
& & N_{F_{q} F_{q}} & & A_{F_{q} W^{\prime}} \\
& & & A_{\nu \nu} & A_{\nu W^{\prime}} \\
A_{F_{1} W^{\prime}}^{T} & \cdots & A_{F_{q} W^{\prime}}^{T} & A_{\nu W^{\prime}}^{T} & A_{W^{\prime} W^{\prime}}
\end{array}\right)\left(\begin{array}{c}
x_{F_{1}} \\
\vdots \\
x_{F_{q}} \\
x_{\nu} \\
x_{W^{\prime}}
\end{array}\right)=\left(\begin{array}{c}
b_{F_{1}} \\
\vdots \\
b_{F_{q}} \\
b_{\nu} \\
b_{W^{\prime}}
\end{array}\right)
$$

Face sets are 2D problems. Only a 5-point stencil needed.

## Comparison of Preconditioners



| $512^{3}$ System | Iterations | Time |
| :---: | :---: | :---: |
| Block Jacobi with Dirichlet-Neumann | 78 | 4977 s |
| Block Jacobi with Neumann-Neumann | 78 | 7882 s |
| Our preconditioner | 10 | 243 s |

## Subdomain Solvers

PC

- A more realistic situation
- FFT in the interior
- PCG for weird boundaries



## Subdomain Solvers

## Memory Issue

- Save the memory as well
- FFT inner solver


## Irregular Subdomain Solvers

- Sparse Cholesky factorization
- PCG

| Degrees of freedom | 2D |  | 3D |  |
| :---: | :---: | :---: | :---: | :---: |
|  | N | Condition Number | N | Condition Number |
| 64 | 8 | 196.829 | 4 | 94.9691 |
| 729 | 27 | 62741.6 | 9 | 1399.99 |
| 4096 | 64 | 79567.4 | 16 | 9696.85 |
| 15625 | 125 | 565988 | 25 | 43741.8 |

## Irregular Subdomain Solvers



## Implementation

- Multi-threading
- Fluid solver
- Domain partitioning
- Solver parameters


## Poisson Solver Test

| $\begin{gathered} \hline \text { Resolution: } 512^{3} \\ \text { Inner solver: ICPCG } \end{gathered}$ | CPU |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
|  | 1 | 8 | 16 | 24 |
| ICPCG | 2846.18 | 1997.04 | 1942.07 | 1934.64 |
| 8*8*8 subdomains | 6166.51 | 954.67 | 548.42 | 449.71 |
| 16*16*16 subdomains | 2607.32 | 417.15 | 259.83 | 230.91 |
| 8*8*8 subdomains(FFT) | 456.75 | 80.40 | 55.82 | 52.43 |
| 16*16*16 subdomains(FFT) | 463.29 | 87.92 | 65.15 | 58.35 |

Machine: $2.5 \mathrm{GHZ}, 24$ core,
2 processor system, 128GB memory.

## Runtime for $512^{3}$



## Runtime for $1024^{3}$



## Parallel Speedup for $1024^{3}$



## Speedup over PCG for $512^{3}$



## Comparison with MGPCG

Table VIII. Comparison of iterations and runtimes between MGPCG and Schur complement solver ( $16^{3}$ subdomains).

|  | MGPCG iter | MGPCG time | Schur iter | Schur time |
| :---: | :---: | :---: | :---: | :---: |
| Scene 1 | 23 | 43.35 s | 15 | 24.56 s |
| Scene 2 | 24 | 44.03 s | 14 | 22.99 s |

The results show moderate performance gains for our method.

We consider that multigrid method has some difficulties in the special treatments required to support thin boundaries.

## Fluid Simulation Test

Table IX. Liquid Simulation Time. Resolution $512^{3}$

| Fluid Scene | MIC0-PCG ParalletT | Schur SerialT | Schur ParallelT | Speedup over MIC0-PCG |
| :---: | :---: | :---: | :---: | :---: |
| Dam break | 261.49 s | 128.59 s | 14.96 s | 17.47 |
| Double dam break | 160.44 s | 126.85 S | 14.51 s | 11.05 |
| Drop objects | 163.63 s | 92.83 s | 11.47 s | 14.26 |
| Obstacles | 371.50 s | 187.66 s | 20.08 s | 18.50 |

Table X. Smoke Simulation Time. Resolution $512 \times 768 \times 512$

| Smoke scene | MIC0-PCG ParallelT | Schur ParallelT | Speedup over MIC0-PCG |
| :---: | :---: | :---: | :---: |
| Plume | 678.99 s | 47.36 s | 14.33 |
| Plume with sphere | 1108.16 s | 49.66 s | 22.31 |

## $512 \times 768 \times 512$


$49.66 \mathrm{sec} / \mathrm{timestep}$

## Conclusion

- We use a domain decomposition approach for greater parallelism
- Create a novel Schur complement preconditioner with high convergence rate
- The use of different linear equation solvers in different flow regions
- High parallelism, low computation time and memory cost


## Limitations

- Does this actually work?
- Seems to. But we haven't proved it.
- We still use PCG, so there's a serial portion.
- Optimal performance for FFT requires powers of 2 DoF
- High overhead for low to mid-res sims


## Still Not Convinced?

- Domain partitioning: efficient implementation
- Domain partitioning: non-uniform subdomains
- Domain partitioning: optimal tile sizes
- Heterogenous computing like Liu et al. 2016
- Distributed computing?
- Lot of data transfer...
- Adaptive grid methods



## 谢谢 <br> Thank you!

